

# A Side Journey to Titan

*Revealing and Breaking NXP's P5x ECDSA Implementation on the Way*

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Journees SIF



# Product Description

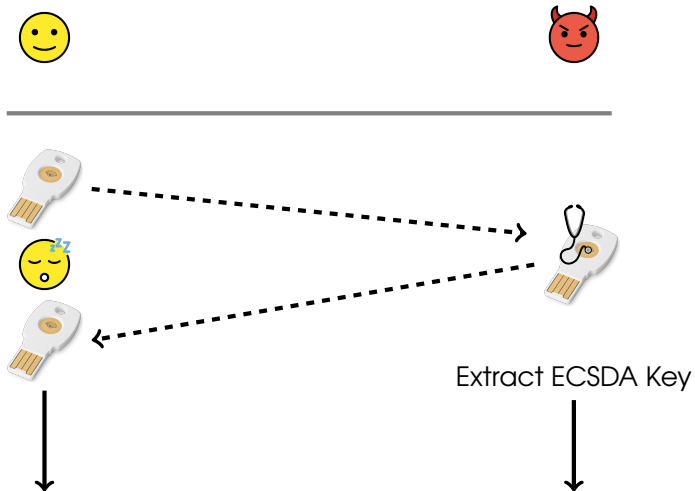
- ▶ **Google Titan Security Key**: hardware FIDO U2F token
- ▶ Hardware token to be used as 2FA for your Google account *and many other services supporting FIDO U2F protocol*
- ▶ 3 versions:
  - ▶ Left: micro-USB, NFC and BLE interfaces
  - ▶ Middle: USB type A and NFC interfaces
  - ▶ Right: USB type C interface



# FIDO U2F Protocol

- ▶ FIDO U2F: open standard for two-factor authentication
  - ▶ Hosted by FIDO alliance (historically developed by Google, Yubico and NXP)
- ▶ FIDO U2F protocol works in two steps - for each account:
  - ▶ *Registration* → ECDSA key pair generation
  - ▶ *Authentication* → ECDSA signature

# Side-Channel Attack Scenario



## Trusted hardware



Titan Security Keys are designed to make the critical cryptographic operations performed by the security key strongly resistant to compromise during the entire device lifecycle, from manufacturing through actual use.

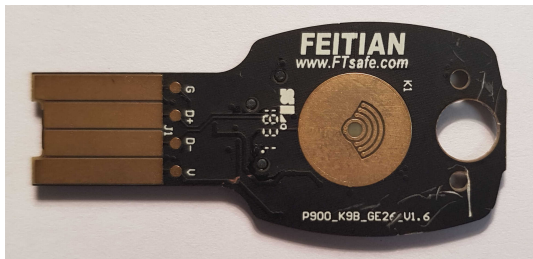
The firmware performing the cryptographic operations has been engineered by Google with security in mind. This firmware is sealed permanently into a secure element hardware chip at production time in the chip production factory. The secure element hardware chip that we use is designed to resist physical attacks aimed at extracting firmware and secret key material.

These permanently-sealed secure element hardware chips are then delivered to the manufacturing line which makes the physical security key device. Thus, the trust in Titan Security Key is anchored in the sealed chip as opposed to any other later step which takes place during device manufacturing.

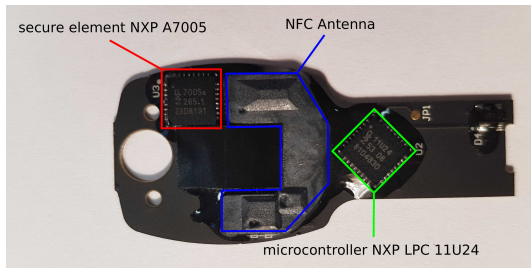
From <https://cloud.google.com/blog/products/identity-security/titan-security-keys-now-available-on-the-google-store>

# Google Titan Security Key Teardown

- ▶ Recto: HW manufacturer is **Feitian**



- ▶ Verso: secure element is **NXP A7005a**



Other FIDO U2F products based on **NXP A700x** chip:

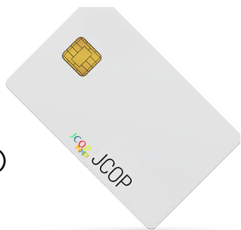
- ▶ Yubico Yubikey Neo
- ▶ Feitian K9, K13, K21, K40

# Similarities with other NXP Products

- ▶ Several NXP JavaCard smartcards (JCOP) can be purchased on the web
- ▶ Those similar to **NXP A700X** are based on **NXP P5x** chips

## NXP J3D081

JCOP 2.4.2 R2  
CC EAL5+ (2015)



- ▶ NXP J3D081\_M59\_DF and variants
- ▶ NXP J3A081 and variants
- ▶ NXP J2E081\_M64 and variants
- ▶ NXP J3D145\_M59 and variants
- ▶ NXP J3D081\_M59 and variants
- ▶ NXP J3E145\_M64 and variants
- ▶ NXP J3E081\_M64\_DF and variants



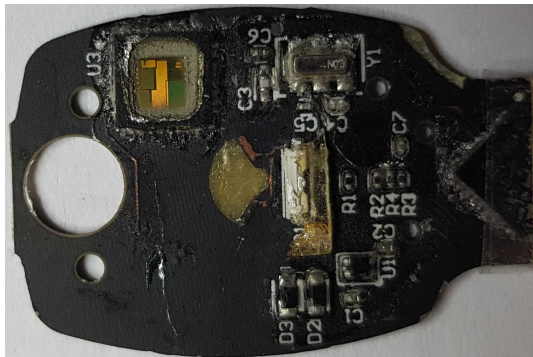
## Rhea

2<sup>nd</sup> largest moon of Saturn after **Titan**

# Titan / Rhea Package Openings

## ▶ *Titan's NXP A700X:*

- ▶ wet chemical opening  
*aluminium tape + fuming nitric acid*
- ▶ **Google Titan Security Key** still alive !



## ▶ *Rhea:*

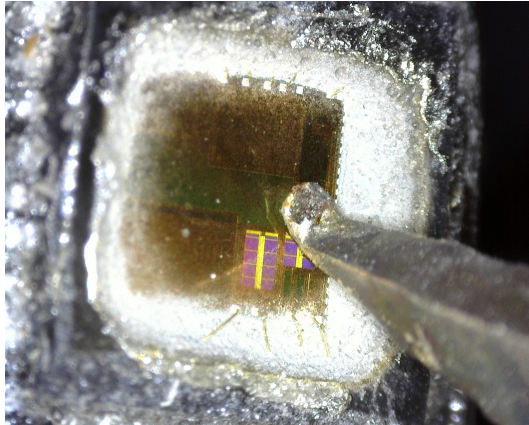
- ▶ mechanical opening  
*scalpel + acetone*
- ▶ **Rhea** still alive !





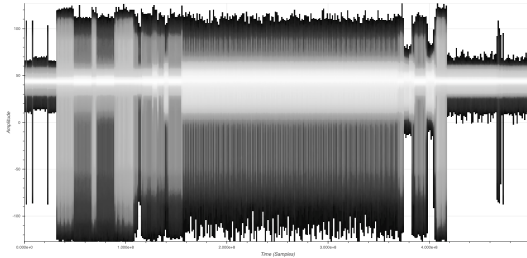
## EM Side-Channel Acquisition Setup (about 10k€)

- ▶ **EM sensor:** Langer ICR HH 500-6 (diam.  $500\mu\text{m}$ , freq. BW 2MHz to 6GHz)
- ▶ **Manual micro-manipulator:** Thorlabs PT3/M 3 axes (X-Y-Z)
- ▶ **Oscilloscope:** PicoScope 6404D, freq. BW 500MHz, SR 5GSa/s

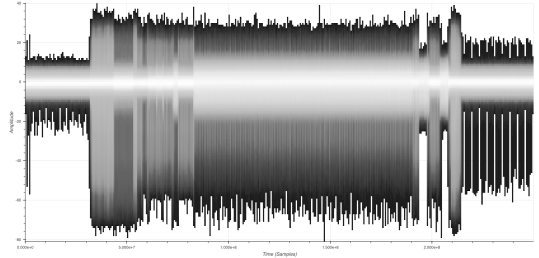


# EM activity of ECDSA signature on Titan / Rhea

## ▶ *Titan*

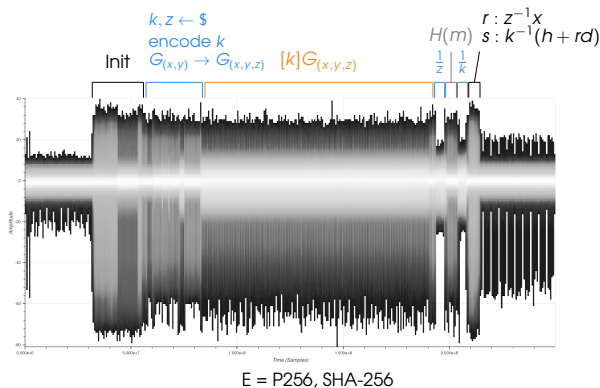


## ▶ *Rhea*



- ▶ ECDSA signature EM activities on *Titan* and *Rhea* look very similar !

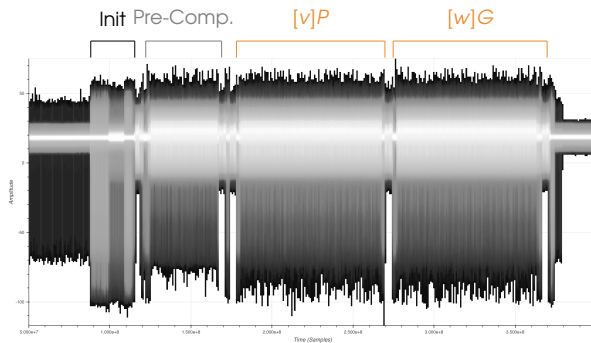
# Rhea – Signature Alg. – EM Radiations



## Scalar Multiplication $[k]G$

- ▶ Constant time algorithm  
*Double-and-Add-Always*
- ▶ 128 iterations for a 256-bit nonce  $k$   
*2 bits of  $k$  by iteration*
- ▶ Randomized point on projective coordinates

# Rhea – Verification Alg. – EM Radiations



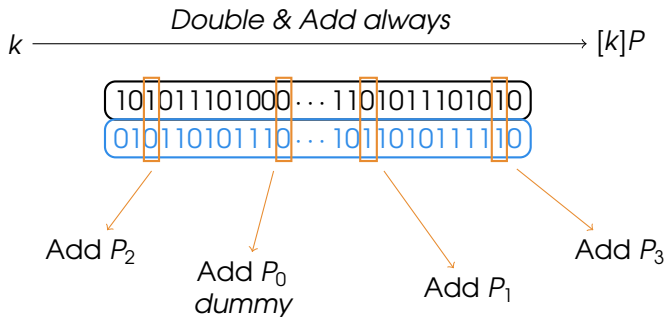
## Scalar Multiplications $[v]P$ and $[w]G$

- ▶ Scalar mult. are **not constant time**  
simple *Double-and-Add*
- ▶ Expensive pre-comp. before  $[v]P$
- ▶ The scalar multiplication algorithm is a **left-to-right comb method** (of width 2)
- ▶ Scalars are not blinded.

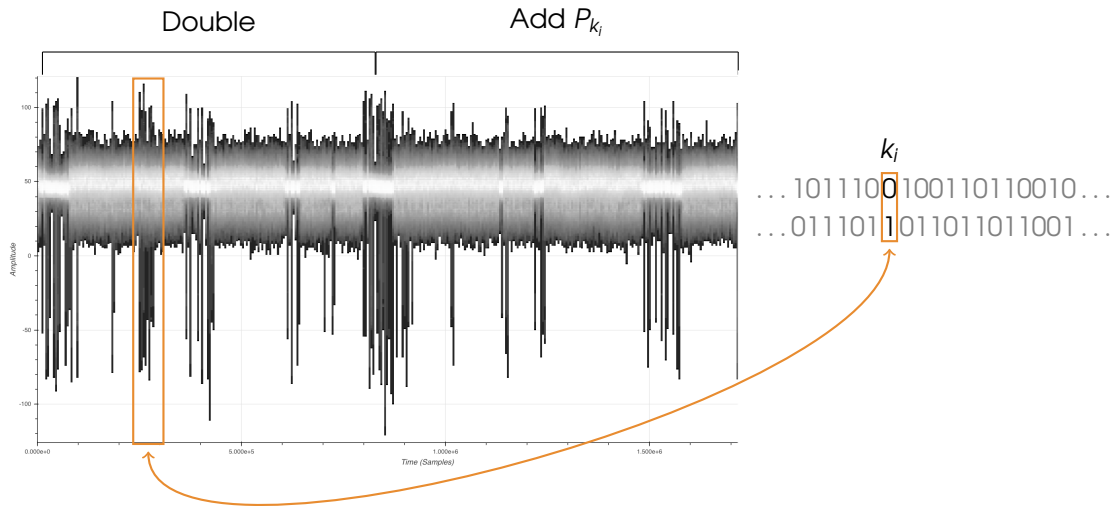
# Scalar Mult. w. Left-To-Right Comb method (width 2)

Precomp:  $P_0, P_1 = P, P_2 = [2^{128}]P, P_3 = P_2 + P_1$

$k =$  101011101000...1101011101010 010110101110...1011010111110

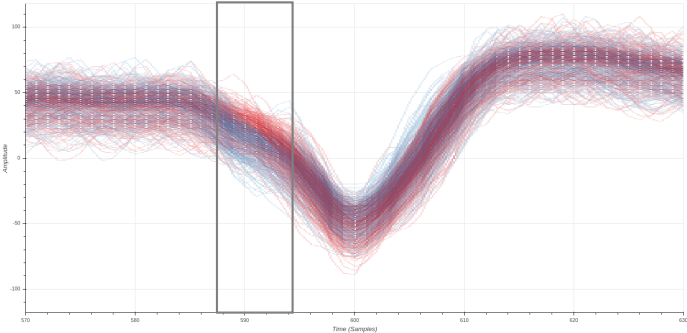


# Rhea – Single Iteration – Leakage Area



# Rhea – Leakage Illustration

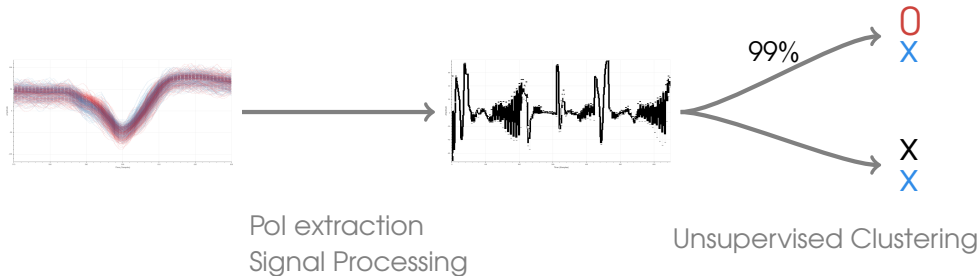
1000 Superposed Iterations – Zoom in Leakage Area



—  $k_i = \frac{0}{X}$

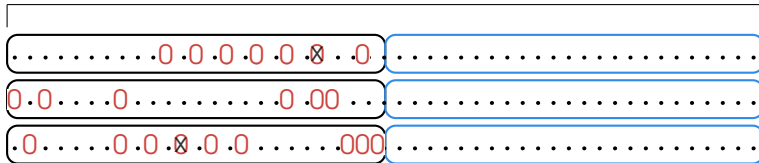
—  $k_i = \frac{1}{X}$

# Single Trace Matching



## 256-bit Nonces

Identified 0 bits  
27.5 by nonce  
in average





# Hidden Number Problem

- ▶ Recovering an ECDSA secret key given some partial knowledge on the nonces can be expressed as a Hidden Number Problem (HNP/EHNP)
- ▶ HNP and EHNP can be defined as games with an oracle
- ▶ The oracle reveals  $x$  and  $f_m(\alpha x)$  for several random values of  $x$   
The player should find the hidden value  $\alpha$
- ▶ HNP:  $f_m$  discloses the  $m$  most significant bits of  $\alpha x$

1101001010.....

- ▶ EHNP:  $f_m$  discloses  $m$  bits of  $\alpha x$ , not necessarily consecutive

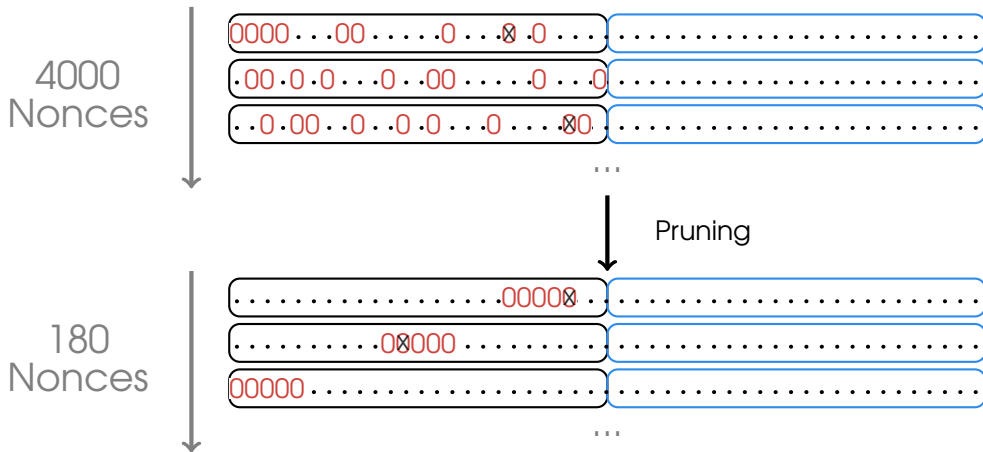
... 1 ... 01 ... 0 ... 101 ... 0 ... 10 ...

# Solving (E)HNP

- ▶ (E)HNP can be reduced to instances of lattice-based problems (SVP, CVP) that may be solved using lattice reduction techniques (LLL, BKZ)
- ▶ # oracle queries and # known bits dictate the size of the lattice and the probability of success
- ▶ In practice EHNP can be solved when the  $m$  bits revealed by the oracle form blocks of sufficiently many consecutive bits

.....0110..... 10110..... 110....

# Rhea – Nonces Selection



# Brute-Forcing the Key

- ▶ LLL reduction (for 80 signatures) takes about 100s
- ▶ 5 errors among 180 available signatures

↔ Brute-force attack on random subsets

## Final Attack

- ▶ Acquisition of 4000 traces:  $\sim 4h$
- ▶ Trace Processing:  $\sim 4h$
- ▶ Brute-force attack:  $\sim 20min$

# Touchdown on Titan

- ▶ Use Rhea parameters for Pol extraction
- ▶ Pruning: from 6000 signatures to 156
- ▶ 7 errors among 156 available signatures

↪ Brute-force attack on random subsets

## Final Attack

- ▶ Acquisition of 6000 traces:  $\sim 6h$
- ▶ Trace Processing:  $\sim 6h$
- ▶ Brute-force attack:  $\sim 30min$

# Impact on Google Titan Security Key

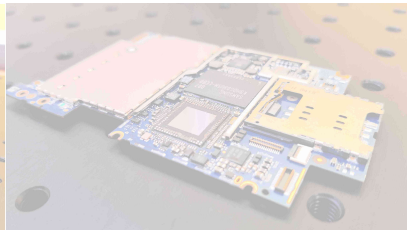
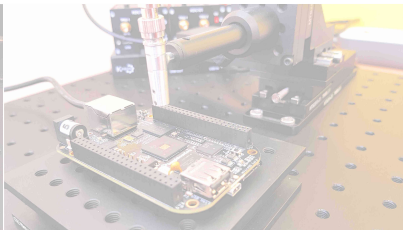
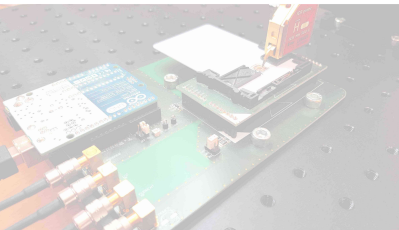
- ▶ Proposed attack allows to physically extract an ECDSA private key linked to application account secured by FIDO U2F
- ▶ But high attack requirements:
  - ▶ Adversary already stolen victim's account login & password
  - ▶ Physical access to **Google Titan Security Key** during 10 hours
  - ▶ Adversary has access to:
    - ▶ Chemical lab
    - ▶ Expensive SCA setup
    - ▶ Custom softwares
- ▶ **So still safer to use Google Titan Security Key than nothing !**

# Attack Mitigations

- ▶ Hardening the NXP P5x cryptographic library:
  - ▶ Blinding of the scalar
  - ▶ Re-randomizing table lookup of precomputed points in comb implementation at each new access
  
- ▶ Use FIDO U2F counter to detect clones:
  - ▶ U2F counter *may* be used for detecting cloned U2F devices
  - ▶ Only limit validity of attack
  - ▶ Dependent on FIDO U2F server implementation

# NinjaLab

*Improve the Security of your Cryptographic Implementation*



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